A SECURED DATA USING CRYPTOGRAPHY IN CLOUD COMPUTING

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Abstract — Cloud computing An innovation which gives the on-request Information Technology administrations for the client through the web. Distributed computing encourages the client by giving the assets of outsider for the sake of framework, equipment and programming over the system. Foundations of Cloud processing makes the client to get to the information anyplace whenever as long as the client's gadget approaches with the web. Such action enhances the utilization of web application which gives "pay as you go" office. Consequently this adaptability makes an effect upon the client and made them to exchange their information to cloud. Be that as it may, it might lay some security issues too. Cryptographic calculations were actualized to beat the security issues and to guarantee the Cloud processing information security. These days numerous procedures of this encryption and unscrambling were proposed to keep up security in cloud information. Here an investigation was made on this cryptographic calculations and a near examination was exhibited.

Index Terms— Cloud Computing, Cryptography, Encryption, Decryption, AES, RSA, MD5.

1 INTRODUCTION

Cloud computing boisterous figuring has created as an Cextraordinarily comprehended technique to encourage broad and voluminous data with the help of shared pool of advantages and huge amassing an area. [1] States that "Distributed computing is another enlisting perspective that depends on virtualization, spread figuring, utility handling and organization arranged building". Advance it is incorporated that distributed computing has created as most basic perspective of the IT business and has pulled in most of the business and the academic network

[2] have portrayed about distributed computing. Distributed computing, no ifs ands or buts, is an extensive term that gives more web benefits. These are confined into three general classes [3]: Infrastructure-as-a-Service , Platform-as-a-Service and Software-as-a-Service. The web is for the most part addressed as the "Cloud". The most section a cloud benefit is used by the clients as and when required, frequently on the hourly introduce. This "on-demand" or "pay as you go" approach impacts the cloud to profit versatile, where end customer can have an inconceivable game plan or unassuming of an adminis-tration the manner in which they need anytime of time and the organization is totally controlled by the provider. Imperative overhauls in each key parts included virtualization passed on enlisting and moreover the improved access to quick web office and what's more weak economy has speeded up the extension of distributed computing altogether.

As cloud figuring appreciates processing as an adequacy, suppliers are building up a common shared gathering of configurable assets, which customers can energetically condition and free as indicated by their changing needs. In this way, both gathering the suppliers and the clients would effortlessly profit by the reuse of figuring assets and diminishing in cost. The cloud benefits that are completed will be executed and tried and true with couple of dangers. Starting advances expected to turn away these dangers. In this way security is the fundamental stress the individuals who need to utilize cloud organizations. As demonstrated by [4] there exist a segment of the crucial security risks that undertaking the Cloud processing usage that spreading spam and malware action of botnets. The other case is the application interfaces that are required to connect with cloud benefits especially that are created by outcasts. These interfaces must outfit the customer with much anchored check, endorsement, encryption and improvement watching frameworks

This paper composed as takes after: Section 2 is the works that identified with distributed computing information security. Area 3 is identified with the distributed computing administrations. Area 4 talks about distributed computing security challenges. Segment 5 clarifies the cryptographic calculations utilized in this examination. Segment 6 delineates the cryptographic calculations execution.

Section 7 shows the conclusion of our work.

2 RELATED WORKS

The most essential objective in [5] is conveying consistent access to control, service, verification and administration organized building organization to end customer. It focused on social occasion the safe and bland plan for that distributed computing stage without knowing its administrations and models. In distributed computing, data is shield from the unapproved singular, disavowal of administration and administration misuse. In [6] the highlights of cloud security systems, assurance issues have focused on specialist co-op side security and proposed the extensible approval tradition for affirmation with RSA computation.

In [7] challenges in surveying the cloud approaches, asset

execution and application outstanding task at hand is portrayed as greatly difficult to achieve, along these lines it proposed, To achieve anchoring and secure access to control, [8] use astoundingly joining methodology of Attributes Based Encryption(ABE), delegate unscrambling and loosened up decoding. It has depicted cryptographic methodology, which give better mystery and security of touchy data outsourced by customer shared on cloud server.

In [9] include every one of security essentials of distributed computing were featured and advising about how to manage the distributed computing security. It have depicted and highlight the general security concern whose made sense of how to comprehend the whole cloud preparing and analyze about the cloud security issues. [10] Have portrayed a security of data to secure information in cloud achieved by Third Party Auditor (TPA), which check the reliability of the dynamic data set away in cloud and play out different looking at assignments in the meantime. Each task on data is annexed with confirmation tag.

In [11] cloud enlisting issues were illustrated i.e. Resolute quality, Availability and Security and it gives the open response for cloud issues. It delineated and depicted all around requested virtualization levels of cloud figuring security. The essential cloud security issues were recognized in [12] and it gives the game plan in cloud handling. It proposes the logical scientific classification design of security and insurance in cloud preparing by segregated the security issue and security plan with accumulated guide. A multi mists database demonstrate has proposed in [13] and it displayed the outline of multi cloud database appear and depicts the layers and sections. [14] have inspected the security issues and discuss all the unquestionable typical for cloud i.e multi-inhabitance, flexibility et cetera and outcast control, by then separate the cloud security necessities i.e. order, respectability and availability ultimately abbreviate the issues in security while cloud handling and cloud plan.3 **Cloud Services**

In the Web, Cloud processing is conclusively giving benefits. The Service models are Infrastructure, Platform and Software is talked about as beneath.

Cloud processing gives various facilitated administrations. The different administration models quickly talked about before have additionally been expounded as beneath, to uncover their hugeness with a scope of security dangers encourage in the overview[15]:

• Infrastructure-as-a-Service (IaaS): It is furthermore specified as Resource Clouds generally give resources and can be scaled up, as organizations to a variety of customers. They essentially supply prevalent virtualization capacities. In this way, extraordinary resources may be offered by methods for an organization line: Data and limit mists convey to the table a time tested access to data of a possibly tremendous size. The accomplishment rate of data gets the opportunity to portray the idea of these cloud servers. As establishment can be logically scaled up or down for the need of usage resources, it gets ready different tenants meanwhile. Also, the benefits that are used are generally charged by the providers.

• Platform-as-a-Service (PaaS): It supplies computation-

al resources by methods for a phase where upon applications and organizations can be urbanized and encouraged. In other way, it supplies all the expected advantages for collect an application and organization through the web, without downloading or presenting it. PaaS generally makes use of over the best APIs to organize the execution of a server encouraging engine which completes and rehashes the execution as shown by buyer requests. As each supplier revealed their own specific API as shown by the individual key potential outcomes, applications delivered for one correct cloud provider can't be enthused to an additional cloud have; there is anyway tries to make more prominent extensive programming models with cloud limits.

• Software-as-a-Service (SaaS): It is also implied as Application or a Service Clouds. SaaS is the model which has the application as a help of its diverse cloud customers by methods for web. The customer utilizes the item out of the case with no compromise or setting up with any system. Organization mists give an execution of unequivocal business limits and business frames as per the essential. These applications are given with unambiguous cloud limits using a cloud structure or stage instead of giving a cloud to them. On and on, sorts of standard application programming value are practical inside a cloud. One most favorable position of SaaS is that it helps in costing less money than truly obtaining the applications to the affiliation.

The three cloud organizations depicted above attract

some significantly basic proportion of threats. This fuses modification of data without ap-propriate support, provoking data cracks or unapproved access to sensitive data. In the event that there ought to emerge an event of genuine data fortification being taken, it is vulnerable that not encoded properly. Unbound access to resources over the cloud may incite unapproved usage of organization,

4 SECURITY CHALLENGES IN CLOUD COMPUTING

Security is the imperative viewpoint for some associations for cloud appropriation. Secrecy, affirmation, respectability, non-renouncement, and openness for client's systems are the general guidelines of security. Get the opportunity to control is another imperative factor for security. There are heaps of security threats to Cloud Service. A single deformity in one client application could empower a threatening software engineer to acquire access for in excess of one client's data. This issue is known as data breaks. The data adversity is another issue that happens when the unapproved customer may delete or change the entire records in the cloud if there is the lack of protection in cloud provider side. Questionable APIs and weak interfaces are another ordinary security challenges in cloud preparing.

Cryptography is likewise a technique for changing over in-arrangement into incoherent frame amid capacity and transmission that it appears to be waste to gatecrasher. The indiscernible data called as figure content. Right when data is gotten by beneficiary, it will appear as unique called as plain content. Changing over to figure content from plain content called encryption and turnaround of this (figure content to plain content) is known as decoding. Encryption occurs at sender's end while decoding occurs at recipient's end. There are three sorts of cryptography calculations[16]. Named Symmetric, Asymmetric and Hashing.

In hashing a signature with fixed length is made with the help hash work or algorithms for the encryption of information. Each message comprises of various hash value, but the hashing has one drawback i.e. once the information is encrypted, it can't be decrypted. This confinement of hashing was evacuated by the algorithm of symmetric and asymmetric. "Secret Key Encryption Algorithm" in symmetric key calculation and single key utilized. i.e. private key, where as in asymmetric algorithms both the keys(Public and Private) are utilized, asymmetric algorithms is otherwise called "Public Key Encryption Algorithm".

5 CRYPTOGRAPHIC ALGORITHMS-COMPARISON

5.1 SYMMETRIC ALGORITHMS

Here Symmetric calculations incorporate a solitary shared mystery key to encode and additionally decode he data and are capable of setting up a lot of information and from preparing standpoint are not to a great degree control escalated, so has cut down overhead on the systems. It has rapid to encode and unscramble the client data with great execution. Symmetric calculations encode the plaintexts as either Stream figures or Block figures with the settled number of 64-bit units.

5.1.1 AES

This Cryptographic algorithm is symmetric block cipher with iterative, which implies that, AES algorithm works by rehashing the same characterized steps again and again. AES algorithm consists with a Secret key. AES calculation takes a shot at a foreordained number of bytes. AES encryption calculation and furthermore a large portion of the encryp-tion calculation is reversible^[17]. With the end goal that, about sim-ilar steps were performed to complete both the scramble and decode in reversible request. The calculation chiefly manages bytes (i.e) it work with bytes, simple to utilize and elucidate. This key is reached out into singular sub keys, which mean a sub keys for all activities. This system is called Key Expansion.

PSEUDO CODE - AES Algorithm

- 1. Choose a password (P) and a salt value(S).
- 2. Get the current time as T.
- 3. Compute key K = S + T.
- 3. Encrypting the password P along with Key K which creates the CT(Cipher Text)
 - CT = AESencrypt(P,K)
- 4. AES encrypt function which does the following process Sub Bytes(SB)-Shift Rows(SR)-Mix Columns(MC) Add Round Key(ARK)
- 5. Decrypt the CT to get plane text Password P by reversing the above process.
- 6. Compute K=S-T
- 7. Plain text password P will obtain by repeating the step 4 in reverse order.
 - P=AESdecrypt(CT,K)

5.1.2 BLOW FISH

Blowfish is one of the Symmetric Cryptographic Algorithm of Block Cipher(BCSCA) and used for en-grave and decode the writings. It utilizes a Variable length key and formed as a speedy and free alternative contrast and existing encryption calculation. Blowfish Algorithm works 16 times. The square size is at first 64 bits then it very well may be stretched out till 448 bits. Each round contains XOR with extension of keys and data encryp-tion[18].

PSEUDO CODE - BlowFish Algorithm

- 1. Input a 64-bit data to Y
- 2. Divide Y into two halves: yL, yR(each 32 bit).
- 3. For Encryption: Compute below step for 16 times starting from P1, P2.....P16 yL = yL XOR PiyR = F(yL) XOR yR
- 4. Swap yL and yR
- 5. After the 16th round, swapping yL and yR again with undo the last swap.
- 6. Compute
 - yR = yR XOR P17 and yL = yL XOR P18.
- 7. Finally, recombine yL and yR to get the cipher text.
- 8. Then getting Decryption same as encryption, but P1 Upto P18 are in the order (reverse).

5.2 ASYMMETRIC ALGORITHMS

Public key cryptography, otherwise called asymmetric cryptography, denotes to a cryptographic algorithm which involves two different keys, one of which is secret key or private key and other one is open key. Despite the fact that disparate, the two areas of this key combina-tion are logically associated. The Public key for en-coding plain substance or to affirm a computerized signature, in like manner the private key is used to decipher the figure message or to make a progressed advanced mark. The expression "Hilter kilter" comes from the usage of different keys to play out these opposite limits each being the converse of the other as showed up diversely in connection to expected "symmetric" cryptography which relies upon a comparative key to perform both.

5.2.1 DIFFIE HELLMAN

Diffie Hellman key trade is a clear system for trading cryptographic keys. This technique grants two client's that have no previous data of one another to commonly set up a typical mystery key over an unverifiable correspondence channel. This key would then have the capacity to be used to encode succeeding correspondences using a symmetric key figure. The calculation is itself limited to the trading of keys[19]. This calculation depends for its feasibility on the inconvenience of figuring discrete logarithms.

PSEUDO CODE - Diffie Hellman Algorithm

- 1. Firstly, S and R are large prime numbers as p1 and p2. These integers kept as secret. S and R can use an insecure channel.
- 2. *S* chooses another random number as large *i.e* (*x*) and calculates c such that
 - $c=p2^x \mod p1$
- 3. S sends the number c to R
- 4. R selects another random integer i.e (y) as independent and find d (i.e)
 - $d=p2^y \mod p1$
- 5. R sends number d to S
- 6. *S* now compute the secrete key Key₁ as follows

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Key₁= d^x mod p1 7. R now computes the secret key Key₂ as follows. Key₂=c ^y mod p1

5.2.2 RSA

RSA is generally used as Public-Key cryptography algorithm defined in 1977. RSA algorithm is employed to encrypt the user information to offer security with the objective that the concerned client can only get the information. First user information is encoded and after that it is deposited in the Cloud. Whenever required, client puts a demand for the information from the Cloud service provider; Cloud supplier verifies and conveys the client data. RSA is also called as block cipher because each message is mapped to a whole number. RSA comprises of Public-Key and Private-Key [20]. In our Cloud atmosphere, all known with public key, while private key known who initially possesses the information. Subsequently, Cloud service provider does the encryption and decryption is handled by the Cloud client or user. Once the information is encoded with the Public-Key, it can be decoded with the equivalent Private-Key only.

PSEUDO CODE - RSA Algorithm

- 1. Choose the prime numbers p and q with distinct
- 2. Calculate the $n = p^*q$.
- 3. Select the *e* as public key that not a factor of which is (p-1) and (q-1)
- Select the public key d which satisfies the (d*e) mod (p-1)*(q-1)=1.
- 5. Encrypting the PT to get CT(Cipher Text) CT=PT e mod n
- 6. Sending Cipher text CT to the receiver.
- 7. Decrypting the CT to get plain text PT CTd mod n

5.3 HASHING ALGORITHMS

Cryptographic Hash capacities are the most fundamental instruments in the field of cryptography and are used to achieve different security targets like validity, Digital Time Stamping, Digital signature, Digital Steganography, pseudo number age et cetera. The hash capacities used in different data preparing applications to achieve diverse security destinations is generously more sweeping than the use of the square figure and the stream figure. Hash limits are to an awesome level of significant and show up in all information security applications. A hash work is a logical strategy that progressions over numerical data into compacted numerical infor-mation. The contribution to the hash work is of selfself-assured length however the yield is constantly of settled length. Characteristics inferred in the hash work additionally called as message process or simply hash esteems.

5.3.1 SHA-3

The Secure Hash Algorithm can be used to make a message known as Message Digest. As decided in that Digital Signature Standard (DSS), the SHA3 algo-rithm joined alongside that Digital Signature Algo-rithm and at whatever point a secured hash calculation is required. The transmitter and expected message of beneficiary in figure and affirm a computerized signature use the SHA3. SHA3 is used for enrolling a data record. Right when a message length under 64 bits of two is input, the SHA3 produces a 160-bit yield known as

Message Digest.

The message process would then have the capacity to be a con-tribution to the DSA, which creates or checks the stamp for the message. Denoting the message procedure rather than the message every now and again improving the adequacy of the technique in that the message procedure is regularly substantially littler in measure than the message. A comparable calculation must be used by a propelled signature as was used by the producer of the mechanized mark [21]. The SHA3 is considered secure in light of the fact that to develop a message which identifies with a given message process, or to find two novel messages which make a comparable message digest. Any change to a message in movement will, with high likelihood, result in an alternate message process, and the mark will neglect to check.

PSEUDO CODE - SHA-3 Algorithm

- 1. Input a Message M, a pointer to the Message p and byte length of M as BL.
- 2. Compute z = 128M + p, $0 \le s \ 128$. If $p \le 111$, the number of calls to update is (M+1) If p > 111, the number of calls to update is (M+2)
- 3. Denote M = floor(x/64) and $s = z \mod 64$, and
- 4. Consider the last block LB as zero
- LB=Null
- 5. Assign the string to the blocks as LB [byte 0] = 0x80 Till LB [byte 15]
- 6. Append(M, LB)
- 7. Compute till M(BL)/128 Update (hash, M) Compute M = M+128
- 8. Now hash will be the Message digest.

5.3.2 MD5

The MD5 calculation creates a 16 byte hash estimation of length 128-piece , which is generally passed on in content arrangement as 32 hexadecimal number digits. Cryptographic applications utilizes MD5 calculation in different ways, and for the most part utilized for confirming information trustworthiness. MD5 calculation forms a variable length to settled length. The message yield will be 128 bits measure. The client message at that point divided into 512 piece squares lumps (i.e) the message extending like 16 times of 32-bit words) so the length can be detachable by 512 piece squares.

padding acts as per the accompanying advances: at first somewhat single as 1, and appended to the end or the last position of the message. This is trailed by as a few quantities of zeros, which is required to get the message length up to 64 bits which is not exactly a different of 512. Whatever remains of the bits with 64 bits and the length of first message, which is modulo of 264. The principal MD5 calculation deals with a 128 piece, apportioned into 32 bit expressions of four. These are set to certain A to D settled constants. The basic calculation at that point hones each Message square of 512 piece to change the state. It includes four comparative stages, as said above, is named as rounds; each round with 16 tasks to view.[22].

PSEUDO CODE - MD5 Algorithm

9. Input the message block *M* of size 512 bits.

10. Split M into 16 32-bit words as M0, M1, M2,... M15.

11. Split the state into four as A,B,C,D

12. Store the state in some variables: $A \rightarrow A', B \rightarrow B', C \rightarrow C'$ and $D \rightarrow D'$

13. Compute the below steps for 64 rounds:

- *i.* Compute T = B + ((A + fi(B, C, D) + Mk + Xi) <<< si).
- *ii.* Rotate the state words: $D \rightarrow A$, $C \rightarrow D$, $B \rightarrow C$, $T \rightarrow B$.

14. Add the stored state values to the state variables:

 $A+A' \to A, B+B' \to B, C+C' \to C, D+D' \to D.$

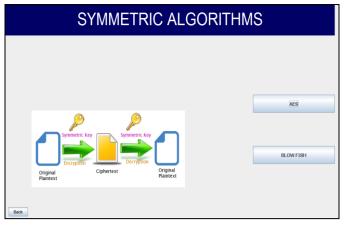
15. Finally that new running state value is the hashed value.

6 EXPERIMENTAL RESULTS

The comparative study of this Cryptographic algorithm was studied and implemented in java environment and experimented the Performance of algorithms(Encryption and Decryption). The evaluation is intended to find the performance of the cryptographic algorithms by dividing the algorithms by their nature as Symmetric Algorithms, Asymmetric Algorithms and Hashing algorithms. The performance calculation for Encryption and Decryption of algorithm was done based on the execution time of each algorithm for different file size.



FIGURE 6.1 - MAIN SCREEN OF THE RESEARCH WORK





| | AES ALGORITHM | | |
|------|--|---|--|
| | Input File : | D FilesiliadrasPetitoner/Service LawG P MITTALActTlissue1(SAW1 M | |
| | Enter the String: | Or Perskamme the basie. Therefore, the awards as wer as the orber to the learnes h to set adde. (5) In the read, the well appears stand allowed and the impugned Awards as a kanned single Judge, date December 14, 2000 are set aside. No cests Cons N.X.M.F. are dicised. | |
| | Key : | MUUDGENER/ITREPOT Encrypt DECRYPT | |
| Back | Text: 0.(qd9603p/9/%0 HomNo6/ b of0-c | DISIVENTIST C CC VISITH CX X 3 No CCC FEG2L_CC2VCS m¥C B°C 19 In the result, the writ appeals stand allowed and the impugned | |

FIGURE 6.3 – ENCRYPTION AND DECRYPTION USING THE AES ALGORITHM

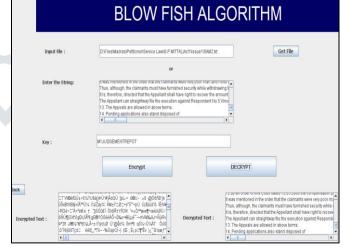
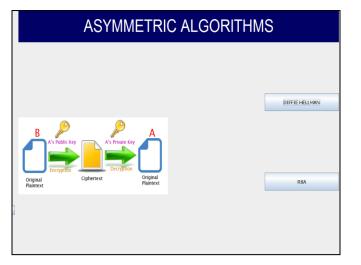


FIGURE 6.4 – ENCRYPTION AND DECRYPTION USING THE BLOWFISH ALGORITHM





| DIFFIE-HELLMAN ALGORITHM | | | | |
|--------------------------|--------------------------------------|------------------------------------|------------------------------|--|
| | Public Keys : | 54 3 KEY EXCHANGE | | |
| | Shared Secret Key : | 27 My Tender Amount is 1,000,00 | | |
| | Enter the Plain Text : | wy remore amount is 1,000,00 | | |
| Encrypted Text : | 0=AF= =>= ===== #)==: #)==: | Decrypted Text : | My Tender Amount is 1,000,00 | |
| Back | ENCRYPT | b | DECRYPT | |

FIGURE 6.6 – ENCRYPTION AND DECRYPTION USING THE DIFFIE HELLMAN ALGORITHM

| The Tender amount is 1,000,00 |
|---|
| 284101110100101114329710911111711011632105115324844848484848484848 |
| 4113-957027-4556-111-67-20-133226-3-62-115115-1-2256600-95-1211062-5-7 4 |
| 28410111010010111432971091111171101163210511532494448484848484848 |
| The Tender amount is 1,000.00 |
| DECRYPT |
| |

FIGURE 6.7 – ENCRYPTION AND DECRYPTION USING THE RSA ALGORITHM

| HASHING FUNCTION | |
|------------------|--------------------|
| Hach fey | <u>8443</u> MD5 |

FIGURE 6.8 - HASHING ALGORITHMS TAKEN FOR STUDY.

| SHA-3 HASHING FUNCTION | | | | |
|------------------------|---|--|--|--|
| Input File : | Get File | | | |
| | or MYTENDER AMOUNTIS 1,000,00 | | | |
| Enter the String: | | | | |
| | HASH | | | |
| Hashed Value : Back | 4538cfs2741e146340495a3380806c2a8c56553554c464276a7743548s1215f | | | |

FIGURE 6.9 – GENERATING HASH VALUE USING THE SHA-3

ALGORITHM

| MD5 HASHING FUNCTION | | | | |
|----------------------|------------------------------------|--|--|--|
| Input File : | Get File | | | |
| Enter the String: | or MY TENDER AMOUNT IS 1,000,00 | | | |
| | HASH | | | |
| Hashed Value : | 2c1a09900833005c0628ce1c4ebc88672 | | | |

FIGURE 6.10 – GENERATING HASH VALUE USING MD5 ALGORITHM

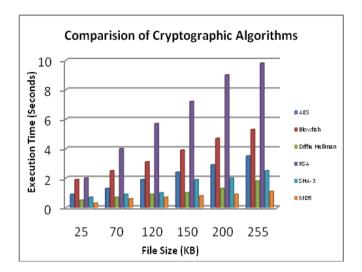


FIGURE 6.11 - CRYPTOGRAPHIC ALGORITHMS-A COMPARISON

7 CONCLUSION

Cryptography is the vital approach of the cutting edge arrange security developments that empower us to send secure data over an inconsistent channel and to guarantee the noteworthy data on the web, extranet, and the intranets. This paper broke down different strategies for data security in the cloud. Dif-ferent encryption methods were proposed by the analysts to make cloud data secure, unprotected were examined. In continuation with that security issues, challenges and besides methods of Encryption Decryption calculations have been made between Symmetric, Asymmetric and Hashing calculations (i.e) AES, Blowfish, Diffie Hellman, RSA, SHA-3 and MD5 figurings to locate the best security calculation for our further procedure as a piece of conveyed processing for making cloud data secure and not to be hacked by assailants.

The calculations of Encryption and Decryption are essential in information security on cloud; here the cryptographic calculations correlation is done in light of estimations of Execution Time. It has been noticed that AES count sets aside the littlest opportunity to execute cloud data. Blowfish and SHA-3 is marginally high in Execution Time, while RSA eats up longest time. The future degree of this work is to find a skilled calculation to impact the data to anchor by merging Diffie Hellman and MD5 figuring and use some pressure calculation for the security of data.

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