

SOLVING FULLY FUZZY ASSIGNMENT PROBLEM USING BRANCH AND BOUND TECHNIQUE

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ABSTRACT

In this paper, we define fuzzy set, fuzzy assignment problem is various classifications of Basic definitions. This paper is organized as follows In section1 we adopt some usual terminology, notations and conventions which will be used later in the section we establish the Branch and Bound Techniques . The fully fuzzy assignment problems is using basic notations of fuzzy numbers and branch and bound techniques have also been presented.

Key words: Fuzzy set, fuzzy Number , Triangular fuzzy number, Basic notations of Assignment problem. .

1.1 INTRODUCTION

The main objective of Assignment problem which is also a special type of linear programming problem is to obtain the optimum assignment for number of tasks (jobs), that is equal to the number of resources (workers) at a minimum cost or maximum profit. This plays an vital role in assigning persons to jobs, classes to rooms, operators to machines etc. An optimal solution to assignment problems, can be achieved using various algorithm such as linear programming, Hungarian algorithm, Neural networks genetic algorithm, etc. In 1965, Zadeh, has introduced fuzzy sets which provided a new mathematical tool to deal with uncertainty of information.

1.2 PRELIMINARIES

1.2.1 Fuzzy set

A fuzzy set is characterized by a membership function mapping the elements of a domain, space, or universe of discourse x to the unit interval $[0,1]$,

$$A = \{ (x, \mu_A(x) : x \in X \},$$

where $\mu_A : X \rightarrow [0,1]$ is a mapping called the degree of membership function of the fuzzy set A . and $\mu_A(x)$ is called the membership value of $x \in X$ in the fuzzy set A . These membership grades are often represented by real numbers ranging from $[0,1]$.

A which maps x to $[0, 1]$.

1.2.2 DEFINITION FUZZY NUMBER

A fuzzy set A defined on the set of real numbers R is said to be a fuzzy number if its membership function $\mu_A : R \rightarrow [0,1]$ has the following characteristics

- i) A is normal. It means that there exists an $x \in R$ such that $\mu_A(x) = 1$
- ii) A is convex. It means that for every $x_1, x_2 \in R$,
- iii) $\mu_A(\lambda x_1 + (1-\lambda)x_2) \geq \min\{\mu_A(x_1), \mu_A(x_2)\}$, $x \in [0,1]$
- iv) μ_A is upper semi – continues
- v) Suppose A is bounded in R .

Arithmetic Operation of Triangular fuzzy number

Let $\tilde{A} = (a_1, a_2, a_3)$ and $\tilde{B} = (b_1, b_2, b_3)$ be two fuzzy triangular fuzzy numbers. then

1. $(a_1, a_2, a_3) \oplus (b_1, b_2, b_3) = (a_1 + b_1, a_2 + b_2, a_3 + b_3)$
2. $k(a_1, a_2, a_3) = (ka_1, ka_2, ka_3)$ for $k \geq 0$
3. $k(a_1, a_2, a_3) = (ka_3, ka_2, ka_1)$ for $k < 0$
4. $(a_1, a_2, a_3) \otimes (b_1, b_2, b_3) = \begin{cases} (a_1 b_1, a_2 b_2, a_3 b_3) & \text{for } a_1 \geq 0 \\ (a_1 b_1, a_1 b_1, a_1 b_1) & \text{for } a_1 < 0, a_3 \geq 0 \\ (a_1 b_3, a_2 b_2, a_3 b_1) & \text{for } a_3 < 0 \end{cases}$

1.3 MATHEMATICAL FORMULATION**1.3.1 The general assignment problem**

Suppose there are 'n' people and 'n' jobs. Each job must be done by exactly one person; also each person can do, at most, one job. The problem is to assign jobs to the people so as to minimize the total cost of completing all of the jobs. The general assignment problem can be mathematically stated as follows:

$$\text{Minimize } Z = \sum_{i=1}^n \sum_{j=1}^n C_{ij} x_{ij}$$

Subject to the constraints

$$\sum_{i=1}^n x_{ij} = 1, \quad \text{for } j = 1, 2, \dots, n.$$

$$\sum_{j=1}^n x_{ij} = 1, \quad \text{for } i = 1, 2, \dots, n$$

where $x_{ij} = 0$ or 1 .

2.4 Fuzzy Assignment Problem

$$\text{Minimize } Z = \sum_{i=1}^n \sum_{j=1}^n C_{ij} x_{ij}$$

Subject to the constraints

$$\sum_{i=1}^n x_{ij} = 1, \quad \text{for } j = 1, 2, \dots, n.$$

$$\sum_{j=1}^n x_{ij} = 1, \quad \text{for } i = 1, 2, \dots, n$$

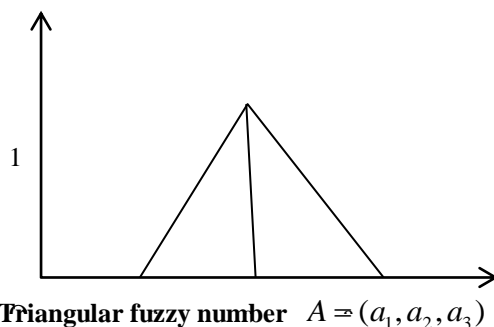
where $x_{ij} = 0$ or 1 .

1.4 Triangular fuzzy number

A fuzzy number A is a triangular fuzzy number denoted by (a_1, a_2, a_3) and its membership function

$\mu_A(x)$ is given below.

$$\mu_A(x) = \begin{cases} \frac{(x-a_1)}{(a_2-a_1)} & \text{if } a_1 \leq x \leq a_2 \\ 1 & \text{if } x = a_2 \\ \frac{(a_3-x)}{(a_3-a_2)} & \text{if } a_2 \leq x \leq a_3 \\ 0 & \text{if otherwise} \end{cases}$$



1. 5. Branching methodology

1. At level k, the column marked as k of the assignment problem, will be assigned with the best row of the assignment problem.
2. If there are tie on the upper bound, then the terminal node at the upper most level is to be considered for further branching.
3. If the maximum upper bound happens to be at any one of the terminal nodes at the (n -1)th level, the optimality is reached. Then the assignments on the path from the root node to that node along with the missing pair of row-column combination will form the optimum solution.

1.6 Numerical Example:

Let us consider a fuzzy assignment problem with rows representing four Machines and column representing four Jobs

Machines/ Job	Job 1	Job 2	Job 3	Job 4
Machine 1	(1,3, 6)	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(8,12,16)	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(13,15,17)	(34,38,40)	(2, 3, 5)	(1, 3, 6)
Machine 4	(27,30,32)	(8,12,16)	(19,22,24)	(19,22,24)

Find an optimal assignment of machines to jobs that will maximize the total profit.(profit in rupees).

Solution:

Consider the fully fuzzy assignment problem and we solve it by branch and bound techniques to get the optimal solution.

Branching

The four different sub problems under the root nodes and upper bounds are shown in fig.1.

$$V_{\epsilon} = \sum_{i,j \in X} c_{i,j} + \sum_{i \in X} \sum_{j \in Y} \max c_{i,j}$$

Where $\epsilon = (1,1)$ $A = (1,1)$ $X = (2,3,4)$ $Y = (2,3,4)$

$$V_{11} = c_{11} + \sum_{i \in (2,3,4)} \sum_{j \in (2,3,4)} \max c_{i,j}$$

Machines/ Job	Job 1	Job 2	Job 3	Job 4
Machine 1	(1,3, 6)	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(8,12,16)	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(13,15,17)	(34,38,40)	(2, 3, 5)	(1, 3, 6)
Machine 4	(27,30,32)	(8,12,16)	(19,22,24)	(19,22,24)

$$P_{11}^1 = c_{11} + \{ \max C_{22} + \max C_{33} + \max C_{44} \}$$

$$= (1, 3, 6) + \{ (27, 30, 32) + (34, 38, 40) + (19, 22, 24) \}$$

Repeated Value has a maximum value

$$P_{11}^1 = (81, 93, 102)$$

Machines/ Job	Job 1	Job 2	Job 3	Job 4
Machine 1	(1,3, 6)	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(8,12,16)	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(13,15,17)	(34,38,40)	(2, 3, 5)	(1, 3, 6)
Machine 4	(27,30,32)	(8,12,16)	(19,22,24)	(19,22,24)

$$P_{21}^1 = c_{21} + \{ \max C_{33} + \max C_{44} + \max C_{11} \}$$

$$= (8, 12, 16) + \{ (34, 38, 40) + (19, 22, 24) + (34, 38, 40) \}$$

$$P_{21}^1 = (95, 110, 120)$$

Machines/ Job	Job 1	Job 2	Job 3	Job 4
Machine 1	(1,3,6)	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(8,12,16)	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(13,15,17)	(34,38,40)	(2, 3, 5)	(1, 3, 6)
Machine 4	(27,30,32)	(8,12,16)	(19,22,24)	(19,22,24)

$$P_{31}^1 = c_{31} + \{ \max C_{44} + \max C_{11} + \max C_{22} \}$$

$$= (13,15,17) + \{ (19, 22, 24) + (34, 38, 40) + (27, 30, 32) \}$$

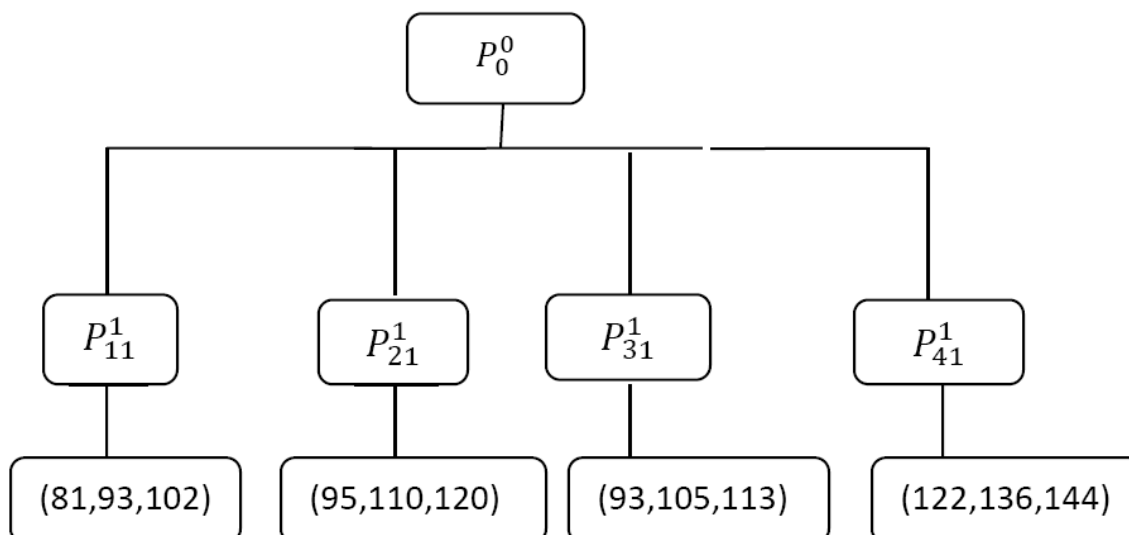
$$P_{31}^1 = (93, 105, 113)$$

Machines/ Job	Job 1	Job 2	Job 3	Job 4
Machine 1	(1,3,6)	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(8,12,16)	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(13,15,17)	(34,38,40)	(2, 3, 5)	(1, 3, 6)
Machine 4	(27,30,32)	(8,12,16)	(19,22,24)	(19,22,24)

$$P_{41}^1 = c_{41} + \{ \max C_{11} + \max C_{22} + \max C_{33} \}$$

$$= (27, 30, 32) + \{ (34, 38, 40) + (27, 30, 32) + (34, 38, 40) \}$$

$$P_{41}^1 = (122, 136, 144)$$



Further branching

Further branching is done from the terminal node which has the greatest upper bound. At this stage, the nodes $P_{11}^1, P_{21}^1, P_{31}^1, P_{41}^1$ are the terminal nodes. The node P_{41}^1 has the greatest upper bound. Eliminate fourth row and first column. Hence further branching from this node is shown as follows.

Machines/ Job	Job 2	Job 3	Job 4
Machine 1	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(34,38,40)	(2, 3, 5)	(1, 3, 6)

$$V_{22} = c_{41} + c_{i2} \sum_{i \in (3,4)} \sum_{j \in (3,4)} \max c_{i,j}$$

Machines/ Job	Job 2	Job 3	Job 4
Machine 1	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(34,38,40)	(2, 3, 5)	(1, 3, 6)

$$P_{12}^2 = c_{41} + c_{12} + \{ \max C_{22} + \max C_{33} \}$$

$$= (27, 30, 32) + \{(8, 12, 16) + (27, 30, 32) + (2, 3, 5)\}$$

$$P_{12}^2 = (64, 75, 85)$$

Machines/ Job	Job 2	Job 3	Job 4
Machine 1	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(34,38,40)	(2, 3, 5)	(1, 3, 6)

$$P_{22}^2 = c_{41} + c_{22} + \{\max C_{33} + \max C_{11}\}$$

$$= (27, 30, 32) + (2, 3, 5) + \{(1, 3, 6) + (34, 38, 40)\}$$

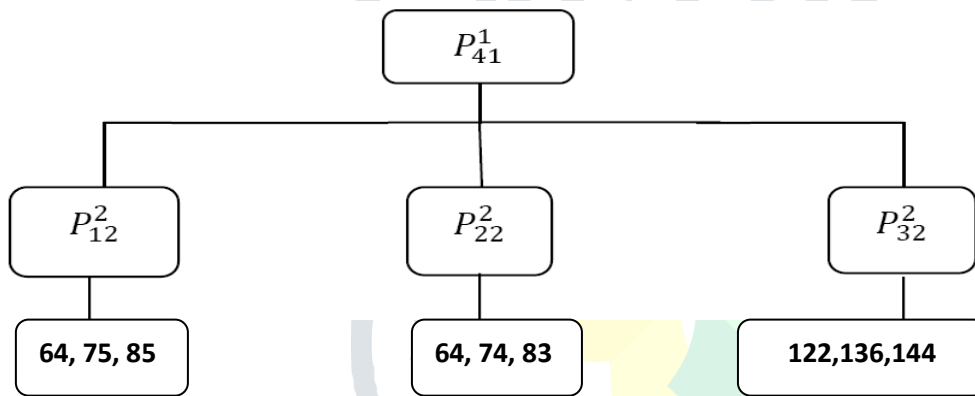
$$P_{22}^2 = (64, 74, 83)$$

Machines/ Job	Job 2	Job 3	Job 4
Machine 1	(8,12, 16)	(20,22, 24)	(34,38, 40)
Machine 2	(2,3,5)	(23,26,28)	(27,30, 32)
Machine 3	(34,38,40)	(2, 3, 5)	(1, 3, 6)

$$P_{32}^2 = c_{41} + c_{32} + \{\max C_{11} + \max C_{22}\}$$

$$= (27, 30, 32) + (34, 38, 40) + \{(34, 38, 40) + (27, 30, 32)\}$$

$$P_{32}^2 = (122, 136, 144)$$



Further branching

Further branching is done from the terminal node which has the greatest upper bound. At this stage, the nodes $P_{41}^1, P_{12}^2, P_{22}^2, P_{32}^2$ are the terminal nodes. The node P_{41}^1 has the greatest upper bound. Eliminate third row and second column.

$$V_{33} = c_{41} + c_{32} + c_{i3} \sum_{i \in (4)} \sum_{j \in (4)} \max c_{i,j}$$

Machines/ Job	Job 3	Job 4
Machine 1	(20,22, 24)	(34,38, 40)
Machine 2	(23,26,28)	(27,30, 32)

Machines/ Job	Job 3	Job 4
Machine 1	(20,22, 24)	(34,38, 40)
Machine 2	(23,26,28)	(27,30, 32)

$$P_{13}^3 = c_{41} + c_{32} + c_{13} + \{\max C_{22}\}$$

$$= (27, 30, 32) + (34, 38, 40) + (20, 22, 24) + \{(27, 30, 32)\}$$

$$P_{13}^3 = (108, 120, 128)$$

Machines/ Job	Job 3	Job 4
Machine 1	(20,22, 24)	(34,38, 40)
Machine 2	(23,26,28)	(27,30, 32)

$$P_{23}^3 = c_{41} + c_{32} + c_{23} + \{ \max C_{11} \}$$

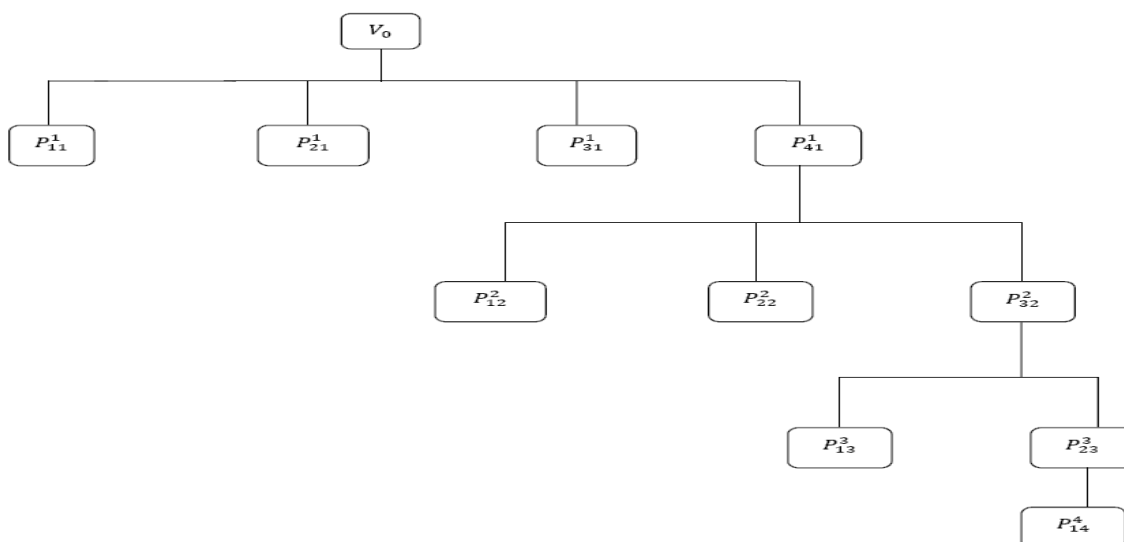
$$= (27, 30, 32) + (34, 38, 40) + (23, 26, 28) + \{ (34, 38, 40) \}$$

$$P_{23}^3 = (118, 132, 140)$$

$$P_{14}^4 = c_{41} + c_{32} + c_{23} + c_{41}$$

$$= (27, 30, 32) + (34, 38, 40) + (23, 26, 28) + (34, 38, 40)$$

$$P_{14}^4 = (118, 132, 140)$$



The optimal assignment

Person 4 → Job1, Person 3 → Job 2, Person 2 → Job 3, Person1 → Job 4.

The optimal cost = (27,30,32) (34,38,40)+(23,26,28)+(34,38,40)
= (118,132,140)

CONCLUSION

In this paper we obtain fuzzy assignment problem was Branch and Bound Method is introduced. This method is effective and easy to understand because of its natural similarity to classical method of solving optimal assignment problems. The method of branch and bound technique is shown in this paper guarantees the correctness and effectiveness of the working produce of the method. The solutions of optimal assignment problems. This method of branch and bound is also easy to apply and understand.

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