# METRO DOMINATION NUMBER OF FRIENDSHIP GRAPH AND CROWN GRAPH

<sup>1</sup>Vishu Kumar M, <sup>2</sup>Nikitha B K, <sup>3</sup>Lakshminarayana S <sup>1</sup>Professor, <sup>2</sup> Research scholar, <sup>3</sup> Assistant Professor Department of Mathematics, School of Applied Science REVA University, Bangalore-560 064, India

Abstract : A Set  $D \subseteq V(G)$  is called dominating set if every vertex  $v \in V - D$  is adjacent to atleast one vertex in D. The dominating set with minimum cardinality is called domination number of G & it is denoted by  $\gamma(G)$ . Let G(V, E) be a graph. A set  $S \subseteq V$  is a resolving set if for every  $u, v \in V(G)$  there exists  $w \in V$  such that  $d(u, w) \neq d(v, w)$ . The resolving set with minimum cardinality is called metric basis and its cardinality is called metric dimension and is denoted by  $\beta(G)$ . A set D which is both resolving set as well as dominating set is called metro dominating set. The minimum cardinality of a metro dominating set is called domination number of G and is denoted by  $\gamma_{\beta}(G)$ . In this paper we determine metro domination of friendship graph, Crown graph and Dragon Graph.

Key Words : Domination, Metric Dimension, Locating Domination, Metro Domination.

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# 1. Introduction

All the graphs considered are simple, finite and connected. A set of vertices *S* resolves a graph *G* if every vertex of *G* is uniquely determined by its vector of distance to the vertices in S. This work undertakes a general study of resolving sets in a friendship graph and Crown Graph. Given a graph G(V, E) and  $u, v \in V(G)$ , d(u, v) denoted the distance between u & v in *G* that is the length of a shortest u - v path. On 1976, Harary and Melter [4] introduce the notation of metric dimension. The vertex set and edge set of a graph *G* denoted by V(G) and E(G). The distance between vertices  $u, v \in V(G)$  is denoted by d(u, v). A vertex  $x \in V(G)$  resolves a pair of vertices  $v, w \in V(G)$  if  $d(v, x) \neq d(w, x)$ . A set of vertices  $S \subseteq V(G)$  resolves *G*, & *S* is a resolving set *S* of *G* with minimum cardinality is a metric basis of *G* & its cardinality is the metric dimension of *G* denoted by  $\beta(G)$ .

# 2. Definitions.

# **2.1 Metric Dimension**

A vertex  $x \in V(G)$  resolves a pair of vertices  $v, w \in V(G)$  if  $d(v, x) \neq g(w, x)$ . A set of vertices  $S \subseteq V(G)$  resolves G, and S is a resolving set S of G with minimum cardinality is a metric basis of G, and its cardinality is the metric dimension of G, denoted by  $\beta(G)$ .

## **2.2 Domination**

Let G(V, E) be a graph. A subset of vertices  $D \subseteq V$  is called a dominating set ( $\gamma$ -set) if every vertex in V - D adjacent to at least one vertex in D. The minimum cardinality of a dominating set is called the domination number of the graph G and is denoted by  $\gamma(G)$ .

#### 2.3 Locating Number

A subset D of V(G) is called a dominating set, if even vertex V - D is adjacent to at least one vertex in D. The minimum cardinality of a dominating set is called the domination number of graph G. The metric dimension of a graph G(V, E) is the cardinality of a minimal subset S of V such that for each pair of vertices u, v of V there is a vertex w in S such that the length of the shortest path from w to u is different from the length of a shortest path from w to v.

The metric dimension of a graph G is also called as a locating number of G and studied its dominating property independently by Slater. A dominating set D is called a locating dominating set or simply LD-set if for each pair of vertices  $u, v \in V - D$ ,  $ND(u) \neq ND(v)$ , where  $ND(u) = N(u) \cap D$ . The minimum cardinality of an LD-set of the graph G is called the locating domination number of G denoted by  $\gamma_L(G)$ .



Figure :1 The graph G and its LD-set



Figure :2 The graph G and its MD-set

#### 2.4 Metro Domination

A dominating set D of V(G) having the property that for each pair of vertices u, v there exists a vertex w in D such that  $d(u, w) \neq d(v, w)$  is called the metro dominating set of G or simply an MD-set. The minimum cardinality of a metro dominating set of G is called metro domination number of G and it is denoted by  $\gamma_{\beta}(G)$ .

#### 2.5 Friendship Graph

A friendship Graph  $F_n$  is a graph.  $F_n$  is a graph that can be constructed by coalescence n copies of the cycle Graph  $C_3$  of length 3 with a common vertex. The friendship graph  $F_n$  is also planar graph with 2n+1 vertices & 3n edges. The vertex set is  $V(F_n) = \{c, v_1, v_{2,...}v_{2n}\}$  and the edge set is  $E(F_n) = \{cv_1, cv_2, cv_3, ... cv_{2n}\} \cup \{v_1v_2, v_3v_4, ..., v_{2i-1}v_{2i}, ..., v_{2n-1}v_{2n}\}$  for  $n \ge 2$  [1]

#### 2.6 Crown graph

An crown,  $C_n \square P_2$  is a graph in which path  $P_2$  is attached at each vertex of the cycle  $C_n$ . Let  $v_0, v_1, v_2, ..., v_{n-1}$  be the vertices of the cycle,  $C_n \square P_2$  such that  $v_i$  is adjacent to  $v_{i+1}$ ,  $0 \le i \le n-2$  and  $v_0$  is adjacent to  $v_{n-1}$ . Also let  $u_0, u_1$  be the vertices of the path which is attached to every vertices of  $C_n$  such that  $u_0$  adjacent to  $u_1$ . We denote the vertices of the path adjacent to the vertex  $v_i$ , as  $u_{i0}, u_{i1}, \ldots, u_{im-1}$ . where  $u_{i0}$  adjacent to  $v_i$ ,  $0 \le i \le n-1$ .

#### 2.7 Dragran Graph

Let  $v_0$   $v_1, v_2, \dots v_{n-1}$ , be the vertices of the cycle of  $D_n m$  such that  $v_i$  is adjacent to  $v_{i+1}$ ,  $0 \le i \le n-2$  and  $v_0$  is adjacent to  $v_{n-1}$ . Also let  $u_0, u_1, u_2, \dots u_{m-1}$  be the vertices of the path or tail of  $D_n m$ . Such that  $u_j$  is adjacent to  $u_{i+1}$ ,  $0 \le j \le m-2$  and  $u_0$  is adjacent to  $v_0$ .

# 2.8 Resolving Set [4]

The subset *S* is a resolving set if  $r\left(\frac{v}{s}\right)$  for every two vertices of *G* have distinct representation. A resolving set of minimum cardinality for *G* is called a minimum resolving set or

a basis for G. If  $S = \{s_1, s_2, ..., s_k\}$  then  $r\left(\frac{v}{s}\right) = (d(v, s_1), d(v, s_2), ..., d(v, s_k)).$ 

### 3. Some known results

In this section we mention some of the known results on metric dimension, domination and metro domination which we use in the subsequent section.

**Theorem 3.1.** (Harary & Melter [4]) The metric dimension of a non trivial complete graph of order n is n-1.

**Theorem 3.2**.(Khuller, Raghavachari, Rosenfeld [8]) The metric dimension of a graph G is 1 if and only if G is a path.

**Theorem 3.3.** (Harary and Melter [4]) The metric dimension of a complete bipartite graph  $K_{mn}$  is m+n-2

**Theorem 3.4.** [9] if G is a graph with no isolated vertices and S is a minimal dominating set of G then V(G)-S is a dominating set of G.

**Theorem 3.5.** (Sooryanarayana and Raghunath [7]) The metro domination number of a graph G is  $\left[\frac{n}{5}\right]$  if and only if G is a cycle.

**Theorem 3.6.** (Sooryanarayana and Raghunath [7]) Let *G* be a graph on *n* vertices. Thus  $\gamma_{\beta}(G) = n-1$  if and only *G* is  $K_n$  or  $K_{1,n-1}$  for  $n \ge 1$ 

**Theorem 3.7.**(Sooryanarayana and Raghunath [7]) For any integer n,  $\gamma_{\beta}(P_4)$  is= $\left[\frac{n}{3}\right]$ .

**Remark 3.8.** For any connected graph  $G, \gamma_{\beta}(G) \ge \max \{\gamma(G), \beta(G)\}$ . [6]

**Theorem 3.9.** For all integer  $n \ge 2$ , dim  $(F_n) = n$ .

## 4. Our results :

**Theorem 4.1** For any integer  $n \ge 3$ ,  $\gamma_{\beta}(F_n) = n$ .

**Proof:** It is obvious that  $\gamma(F_n) = 1$ . (Central vertex 'c' dominates all the vertices of  $F_n$ ). Also by theorem 3.9 dim $(F_n) = n$ . By remark 3.8.

$$\begin{aligned} \gamma_{\beta}(F_n) &\geq \max(\beta, (F_n), \gamma(F_n)) \\ \gamma_{\beta}(F_n) &\geq \max(n, 1). \\ \gamma_{\beta}(F_n) &\geq n. \end{aligned} \tag{1}$$

To prove the reserve inequality we define a dominating set.  $D = \{v_{2k-1} | k \ge 1, 1 \le k \le 2n-1\}$  he above set serves as a domination set, which is also a resolving set by theorem 3.9.

 $\therefore \gamma_{\beta}(F_n) \le n$ 

Thus from (1) & (2)  $\gamma_{\beta}(F_n) = n$ .

**Lemma 4.2** For any integer  $n \ge 3$ ,  $\gamma(C_n \Box P_2) = n$ .

**Proof:** Since the crown graph  $C_n \square P_2$  contains *n*-pendent vertices, to dominates these pendent vertices we require minimum *n* vertices of  $C_n \square P_2$  thus  $\gamma(C_n \square P_2) \ge n$ . (1)

(2)

To prove the reverse inequality we define a set  $D = \{ V_k / 0 \le k \le n-1 \}$ . The above set serves as a dominating set of  $C_n \Box P_2$  thus  $\gamma(C_n \Box P_2) \le n$ . (2) From (1) and (2)  $\gamma(C_n \Box P_2) = n$ .

**Theorem 4.3.** For any integer  $n \ge 3$ ,  $\beta(C_n \Box P_2) = 3$  **Proof : Case-1 for odd** nWe choose a subset  $S = \{u_{00}, u_{10}, u_{\lfloor \frac{n}{2} \rfloor 0}\}$ 

and we must show that dim $(C_n \Box P_2) = 3$ 

For  $n \ge 3$ . By definition 2.7, we got the representations of vertices of graph  $C_n \square P_2$  with respect to *S* are  $r(v_0 \setminus s) = (1, 2, \lfloor \frac{n}{2} \rfloor)$ 

$$r(v_{1} \setminus s) = (2,1, \left\lfloor \frac{n}{2} \right\rfloor)$$

$$r(v_{1} \setminus s) = (2,1, \left\lfloor \frac{n}{2} \right\rfloor)$$

$$r(v_{i} \setminus s) = (i+1, i, \left\lfloor \frac{n}{2} \right\rfloor - (i-1))$$

$$2 \le i \le \lfloor n/2 \rfloor$$

$$r\left(v_{\left\lfloor \frac{n}{2} \right\rfloor} \setminus s\right) = \left(\left\lfloor \frac{n}{2} \right\rfloor, \left\lfloor \frac{n}{2} \right\rfloor, 2\right)$$

$$r(v_{i} \setminus s) = (2\left\lfloor \frac{n}{2} \right\rfloor - i, 2\left\lfloor \frac{n}{2} \right\rfloor - (i-1), i - (\left\lfloor \frac{n}{2} \right\rfloor - 1))$$

$$\left[\frac{n}{2} \right\rfloor + 1 \le i \le n-1.$$

$$r(u_{00} \setminus s) = (4,3, \left\lfloor \frac{n}{2} \right\rfloor)$$

**Case – 2 for even** n

We choose a set  $S = \left\{ u_{00}, u_{10}, u_{\left(\frac{n-2}{2}\right)0} \right\}$ By definition 2.7, we got the representations of vertices of graph  $C_n \Box P_2$  with respect to S are  $r(v_0 \setminus s) = (1, 2, \frac{n}{2})$   $r(v_1 \setminus s) = (2, 1, \frac{n-2}{2})$   $r(v_1 \setminus s) = (i + 1, i, \frac{n}{2} - i)$   $2 \le i \le \frac{n-2}{2}$   $r(v_n \setminus s) = (\frac{n+2}{2}, \frac{n}{2}, 2)$  $r(v_i \setminus s) = (n - (i - 1), n - (i - 2), i - \left(\frac{n-4}{2}\right))$   $\frac{n+2}{2} \le i \le n - 1$ 

From the above (case 1 and case 2), the representations of vertices of cycle  $C_n$  in the graph  $C_n \Box P_2$  are distinct. Also we can observe that the representation of the vertices  $u_{i0} (0 \le i \le n-1)$  are also distinct. This implies that *S* is resolving set, but it is not necessarily the lower bound. Thus the upper bound is dim  $(C_n \Box P_2) \le 3$ .

Now we show that  $\dim(C_n \Box P_2) \ge 3$ .

(1)

(2)

Let  $S = \{u_{00}, u_{10}, u_{\left(\frac{n-2}{2}\right)0}\}$ . Is a resolving set which is |S| = 3. Assume that  $S_1$  is another minimum resolving set or we can denote  $|S_1| < 3$ . If we choose an ordered set  $S_1 \subseteq S - \{u_{i0}\}, i = 0, 1, \left(\frac{n-2}{2}\right)$  so that there are two vertices which are of same representation.  $S_1$  is not a resolving set is a contradiction with assumptions.

Thus the lower bond is  $\dim(C_n \Box P_2) \ge 3$ .

From above proving we conclude that  $\dim(C_n \Box P_2) = 3$ .

**Theorem 4.4.** For any integer  $n \ge 3$ ,  $\gamma_{\beta}(C_n \Box P_2) = n$ . **Proof:** By lemma 4.3 dim $(C_n \Box P_2) = 3$ , also by lemma 4.2  $\gamma(C_n \Box P_2) = n$ By remark 3.8  $\gamma_{\beta}(C_n \Box P_2) \ge \max\{\beta(C_n \Box P_2), \gamma(C_n \Box P_2)\}.$ 

$$\geq \max \{3, n\}$$

$$\gamma_{\beta}(C_n \Box P_2) \ge n$$

To prove the reverse inequality we defined a dominating set  $D = \{\frac{u_{i0}}{0} \le i \le n-1\}$ .

We note that the above set serves as a dominating set which also a resolving set  $\gamma_{\beta}(C_n \Box P_2) \le n$ 

Thus  $\gamma_{\beta}(C_n \Box P_2) = n.$ 

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