

“Development of Matrix for Cryptography”

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Abstract:

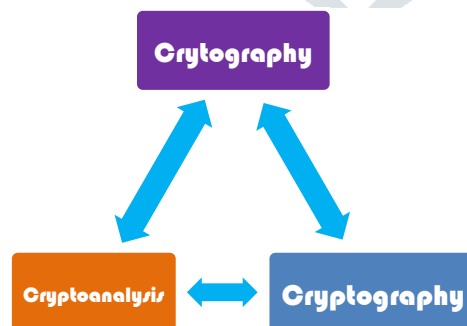
In this article, we have planned to use linear algebra/Matrix for solving cryptographic algorithm, an capable data encryption and data decryption algorithm to protect the communication with the help of key passed between dispatcher and recipient. In this article we are using Hill cipher method to tackle encryption and decryption data.

Keywords: cryptography, Plaintext, cipher text, Encryption, Decryption

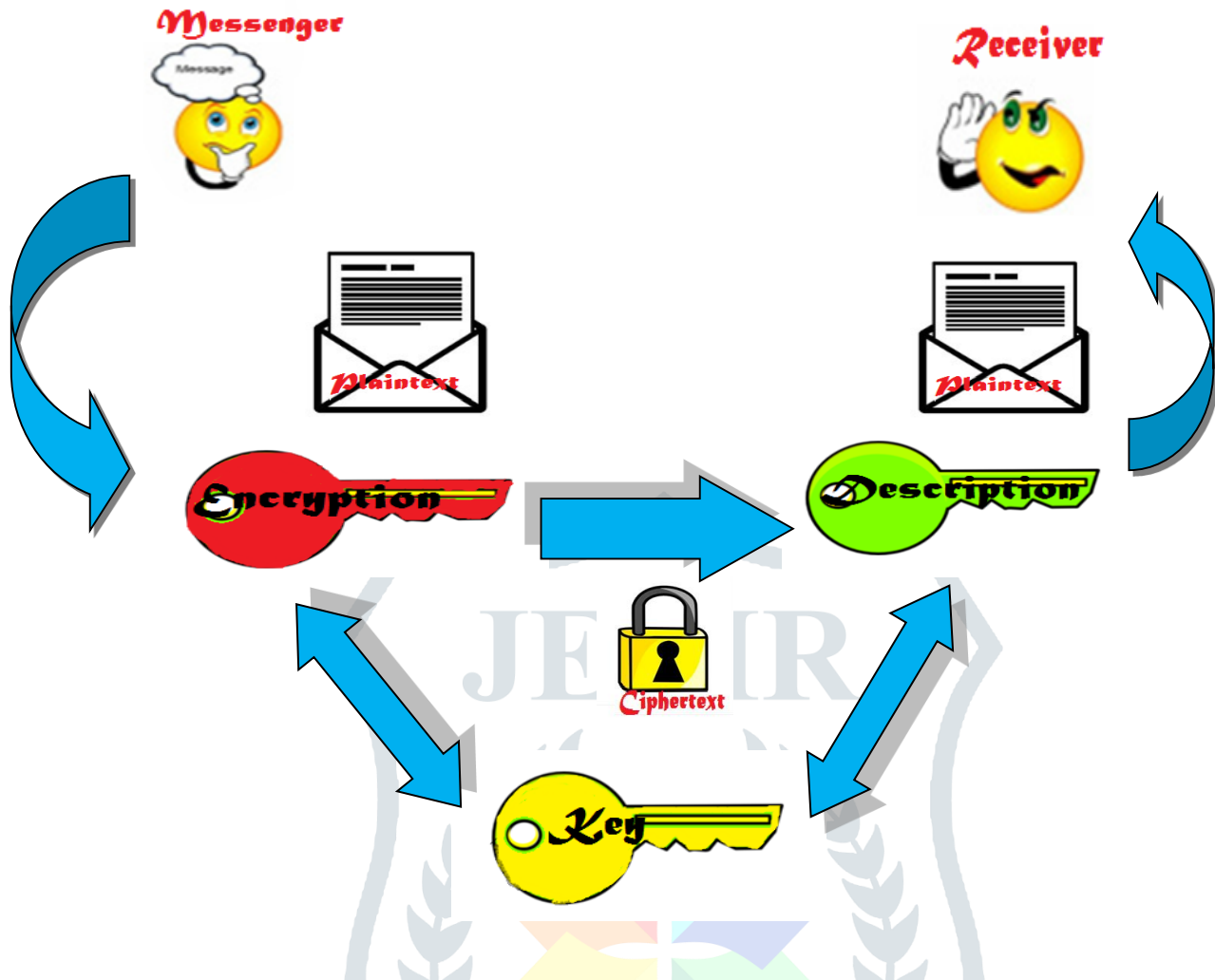
I INTRODUCTION: Cryptography originated about 4000 year ago. Cryptography an Ancient Greek translit “Kryptos” means “Hidden, secret” and graphein means “to write”⁵. Today ‘Cryptography’ is universal in our lives without most of us realizing it. The essential aspect of ‘Cryptography’ has remained same all the way through time which is to hide information in transit and make it presented only for the intended recipients¹. The **encoding** of a note is the invention of the message. It is a system of coded meanings, and in order to create that, the sender needs to understand how the world is comprehensible to the members of the listeners². The **decoding** of a note is how an listeners member is able to understand, and understand the message. It is a process of understanding and conversion of coded information into a understandable form. The messages are called **Plaintext**³.

The messages after coading are called **cipher text**. The process of converting from plain text to coading message is called **enciphering(encoding)**.it is reverse process of getting plain text from cipher text is called **deciphering(decoding)**⁴.

- **Cryptography:** procedure of making and using codes to secure broadcast of information⁶
- **Encryption:** converting original message into a form unreadable by unauthorized individuals⁶
- **Cryptanalysis:** procedure of obtaining original message from encrypted message without knowing algorithms or keys⁶
- **Cryptology:** art of encryption; combines cryptography and cryptanalysis ethics of Information Security⁶



It is the practice and study of hiding information. In this type of Encryption, a pair of keys known as Public Key and Private Key is used. As name indicates, Public Key is shared and identified to everyone where as Private Key is With the person himself^{7,6}.



II OBJECTIVES OF CRYPTOGRAPHY:

Confidentiality: the information cannot be known by anyone for whom it was unexpected⁵.

Integrity: the information cannot be altered in storage or transfer between dispatcher and intended receiver without the modification being detected⁵

Non-repudiation: the originator of the information cannot reject at later stage his or her intentions in the making or broadcast of the information⁵

Authentication: the dispatcher and recipient can verify each other's identity and source/destination of the information.

All are needed in various applications such as Military communications, Radio communication, telephonic communication, Network communication, Mobile communication and internet^{5,6}.

III RESEARCH METHODOLOGY

Hill cipher is a polygraphic substitution cipher that is based on linear algebra. It is the first polygraphic cipher that's able to work on more than three symbols at once^{5,13}

3.1 A history of hill cipher

It was invented by Lester S. Hill in 1929. Hill cipher is an example of block cipher that acts on groups of letters¹¹. It can work on any sized block, including digraphs & triagraphs. The ability to extendable to work with different sized blocks of letters is main advantage of this cipher^{5,6}.

3.2 Method to encrypt using Hill cipher

To encrypt you must change your keyword into a key matrix, 2 by 2 for digraphs or 3 by 3 for trigraphs¹². Then turn the plaintext into digraphs or trigraphs & each of these into column vector. Then perform matrix multiplication modulo the length of alphabet on each vector. Then finally convert these vectors back into letters to produce the ciphertext^{5,6}.

Encryption¹⁰

Cipher Text = (plain Text X key) mod 26

E.g. C = AB(mod 26)

Description⁹

Plain Text = (Cipher Text X Key⁻¹) mod 26

E.g. B = A⁻¹ C (mod 26)

3.3 Example of hill cipher encryption using 3 by 3 matrix⁸**Step I**

Encryption

.Message: ATTACK IS TONIGHT

$$\text{Key} = \begin{bmatrix} 3 & 10 & 20 \\ 20 & 9 & 17 \\ 9 & 4 & 17 \end{bmatrix}$$

Step 2

Encryption

Message:: ATTACK IS TONIGHT

Assign: A-Z 0-25

$$\begin{bmatrix} A & T & T \\ A & C & K \\ I & S & T \\ O & N & I \\ G & H & T \end{bmatrix} = \begin{bmatrix} 0 & 19 & 19 \\ 0 & 2 & 10 \\ 8 & 18 & 19 \\ 14 & 13 & 8 \\ 6 & 7 & 19 \end{bmatrix}$$

Step 3

.....Encryption

Message: ATTACK IS TONIGHT

Cipher Text = (Plain Text X key) Mod 26

$$= \begin{bmatrix} 0 & 19 & 19 \\ 0 & 2 & 10 \\ 8 & 18 & 19 \\ 14 & 13 & 8 \\ 6 & 7 & 19 \end{bmatrix} \times \begin{bmatrix} 3 & 10 & 20 \\ 20 & 9 & 17 \\ 9 & 4 & 17 \end{bmatrix} \text{ mod } 26$$

Step 4

.....Encryption

$$= \begin{bmatrix} 0 & 19 & 19 \\ 0 & 2 & 10 \\ 8 & 18 & 19 \\ 14 & 13 & 8 \\ 6 & 7 & 19 \end{bmatrix} \times \begin{bmatrix} 3 & 10 & 20 \\ 20 & 9 & 17 \\ 9 & 4 & 17 \end{bmatrix} \text{ mod } 26$$



$$C_{11} = 0*3 + 19 * 20 + 19*9 = 551 \text{ Mod } 26 = 05$$

$$C_{12} = 0* 10 + 19*9 + 19 * 4 = 247 \text{ mod } 26 = 13$$

$$C_{13} = 0 *20 + 19*17 + 19*17 = 646 \text{ mod } 26 =22$$

$$C_{21} = 0*3 + 2*20+10*9=130 \text{ Mod } 26=0$$

$$C_{22}=0*10+2*9+10*4=58 \text{ Mod } 26=6$$

$$C_{23}=0*20+2*17+10*17=204 \text{ mod } 26=22$$

$$C_{31} = 8*3 + 18*20+19*9=555\text{mod}26=9$$

$$C_{32}=8*10+18*9+19*4=318 \text{ mod}26=6$$

$$C_{33}=8*20+18*17+19*17=789 \text{ mod } 26=9$$

$$C_{41}=14*3+13*20+8*9=374 \text{ mod } 26=10$$

$$C_{42}=14*10+13*9+8*4=289 \text{ mod } 26 =3$$

$$C_{43}=14*20+13*17+8*17=697 \text{ mod } 26 = 13$$

$$C_{51}=6*3+7*20+19*9=329 \text{ mod } 26 = 17$$

$$C_{52}=6*10 +7*9+19*4=199 \text{ mod } 26= 17$$

$$C_{53} = 6*20+7*17+19*17= 562 \text{ mod } 16=16$$

ATTACK IS TONIGHT --> FNDAGW JG JKDNRRQ

$$\text{Cipher Text} = \begin{bmatrix} 5 & 13 & 22 \\ 0 & 6 & 22 \\ 9 & 6 & 9 \\ 10 & 3 & 13 \\ 17 & 17 & 16 \end{bmatrix}$$

Decryption

$$\text{Key} = \begin{bmatrix} 3 & 10 & 20 \\ 20 & 9 & 17 \\ 9 & 4 & 17 \end{bmatrix}$$

$$\text{Det Key} = \begin{vmatrix} 3 & 10 & 20 \\ 20 & 9 & 17 \\ 9 & 4 & 17 \end{vmatrix} \text{ mod } 26 = -1635 \text{ mod } 26 = -23 \text{ mod } 26 = 03$$

$$\{\text{Det (key)}\}^{-1} = (03)^{-1} \text{ mod } 26 = 09$$

$$\text{Transpose of key} = \begin{bmatrix} 3 & 20 & 9 \\ 10 & 9 & 4 \\ 20 & 17 & 17 \end{bmatrix}$$

$$\text{Minor of Transpose of key} = \begin{bmatrix} 85 & 90 & -10 \\ 187 & -129 & -349 \\ -1 & -78 & -173 \end{bmatrix}$$

$$\text{Cofactor Matrix of Transpose of key} = \begin{bmatrix} 85 & -90 & -10 \\ -187 & -129 & 349 \\ -1 & 78 & -173 \end{bmatrix}$$

$$\begin{aligned}
 (\text{key})^{-1} &= (\text{Det Key})^{-1} \times \text{Adj Key} = \left\{ 09 * \begin{bmatrix} 85 & -90 & -10 \\ -187 & -129 & 349 \\ -1 & 78 & -173 \end{bmatrix} \right\} \text{mod } 26 \\
 &= \begin{bmatrix} 765 & -810 & -90 \\ -1683 & -1161 & 3141 \\ -9 & 702 & -1557 \end{bmatrix} \text{mod } 26 = \begin{bmatrix} 11 & 22 & 14 \\ 7 & 9 & 21 \\ 17 & 0 & 3 \end{bmatrix}
 \end{aligned}$$

Plain text = (Cipher Text X key⁻¹) mod 26

$$\begin{aligned}
 &= \left\{ \begin{bmatrix} 5 & 13 & 22 \\ 0 & 6 & 22 \\ 9 & 6 & 9 \\ 10 & 3 & 13 \\ 17 & 17 & 16 \end{bmatrix} * \begin{bmatrix} 11 & 22 & 14 \\ 7 & 9 & 21 \\ 17 & 0 & 3 \end{bmatrix} \right\} \text{mod } 26 \\
 &= \begin{bmatrix} 520 & 227 & 409 \\ 416 & 54 & 192 \\ 294 & 252 & 279 \\ 352 & 247 & 242 \\ 578 & 527 & 643 \end{bmatrix} \text{mod } 26 = \begin{bmatrix} 0 & 19 & 19 \\ 0 & 2 & 10 \\ 8 & 18 & 19 \\ 14 & 13 & 8 \\ 6 & 7 & 19 \end{bmatrix} = \text{ATTACK IS TONIGHT}^5
 \end{aligned}$$

Likewise we tackle encryption and decryption of our cryptography examples. This methodology is very helpful to solve crypto problems.

IV. Applications of Cryptography:

- **Secure communication:**

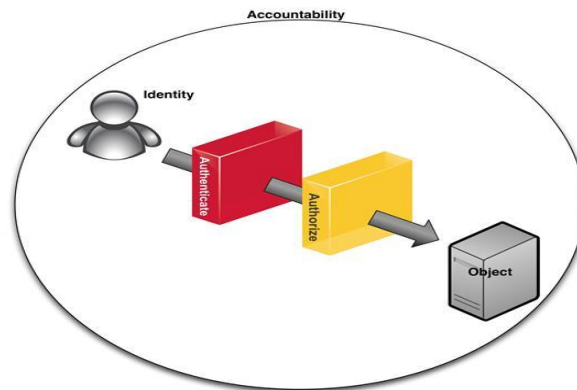
To stop listen in-war time communication and commerce transaction.



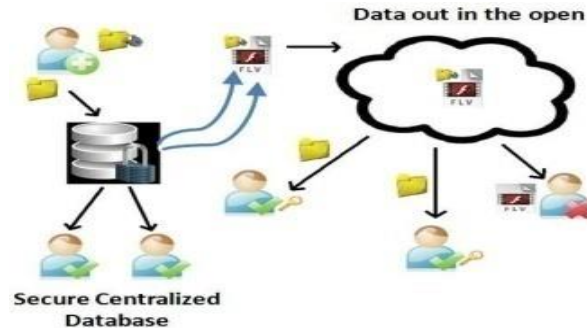
Military communications⁵.

- **Identification & Authentication:**

Checking the integrity⁵



- **Secret sharing/ data hiding:**
Secrete something that has been written⁵.



- **E-commerce/ E-payment:**



- **Certification:**

Certification is a system by which trusted agents such as certifying authorities assurance for Unidentified agents, such as users⁵.

V Conclusion:

In this article, we have planned to use linear algebra/Matrix for solving cryptographic algorithm, an capable data encryption and data decryption algorithm to protect the communication with the help of key passed between dispatcher and recipient. Also communication with any number of words having any digit of character can is encrypts and decrypt by dispatcher and recipient. With data encryption, data owner can use the benefits of communication splitting to digit of words such that to decrease storage and computational overheads. The hill cipher methodology can be used to tackle crypto problem, using matrices it is possible to solve encryption and decryption problem.

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